CS 130 Homework 4

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Practice Problems

You are not required to turn these in.

Gersting, 6e: Section 1.2, Exercises 12–22 and Section 1.4, Exercises 10–27

Book of Proof: N/A (this book is light on formal logic)

Note that Gersting uses a different syntax than we've been using in class. Basically, they write $A \wedge B \wedge C \longrightarrow D$ to mean what we've written as A, B, $C \vdash D$. They're functionally equivalent (cf. the Deduction Theorem and Exportation), but may be hard to read at first—especially because the textbook also elides parentheses and square brackets a lot.

Turn-in Problems

Prove the following theorems using the predicate logic proof system. You may use any previously proven result from the homework/lecture and any theorem from the list provided at http://www.csupomona.edu/~ajvondrak/cs/130/12/winter/theorems.pdf.

$$\begin{array}{cccc} 1 & \vdash & (\forall x [\forall y [A(x,y)]] \Longrightarrow \forall y [\forall x [A(y,x)]]) \\ \hline 2 & (a) \neg \exists x [A(x)] & \vdash & \forall x [\neg A(x)] \\ & (b) \forall x [\neg A(x)] & \vdash & \neg \exists x [A(x)] \\ & (c) \exists x [\neg A(x)] & \vdash & \neg \forall x [A(x)] \\ & (d) \neg \forall x [A(x)] & \vdash & \exists x [\neg A(x)] \\ \hline 3 & \vdash & (\exists x [(A(x) \Longrightarrow B(x))] \Longrightarrow (\forall x [A(x)] \Longrightarrow \exists x [B(x)])) \\ \hline 4 & \vdash & \forall y [\exists x [(A(y,x) \Longrightarrow A(y,y))]] \end{array}$$

Hints:

- For $\lfloor 2 \rfloor$ (a)-(d), instead of writing down and/or looking for a \exists in your formula, use the "expanded" form.
- For 2(c), prove the lemma

$$\forall x[A(x)] \qquad \vdash \qquad \forall x[\neg \neg A(x)]$$

then apply the Deduction Theorem to get the tautology

$$\vdash \qquad (\forall x[A(x)] \implies \forall x[\neg \neg A(x)])$$

which you can use as a line in your proof.

• For 2 (d), prove the lemma

$$\forall x[\neg \neg A(x)] \qquad \vdash \qquad \forall x[A(x)]$$

then apply the Deduction Theorem to get the tautology

$$\vdash \qquad (\forall x[\neg \neg A(x)] \implies \forall x[A(x)])$$

which you can use as a line in your proof.

• In general, instead of proving something of the form $\vdash (A \implies B)$ directly, it's usually easier to prove $A \vdash B$ so you actually have a hypothesis. Then use the Deduction Theorem on this result to show $\vdash (A \implies B)$ must be true.