# Recursion CS 130

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#### Recursion

#### Definition

Recursion is the process of defining something in terms of itself

- The base case is the simplest instance of the definition, which requires no self-reference
- The recursive case is a more complex instance of the definition, which relies on self-reference to a simpler case (i.e., an instance closer to being the base case)

#### Example (Recursive Exponentiation)

Suppose we're dealing with natural numbers (0, 1, 2, ...). Exponentiation can be defined recursively upon the operands *a* and *b*:

$$a^0 = 1$$
 (Base Case:  $b = 0$ )  
 $a^b = a \times a^{b-1}$  (Recursive Case:  $b > 0$ )

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Using the recursive definition of exponentiation, what is the value of the following highlighted subexpression?

$$2^3 = ?$$

(A) 8  
(B) 
$$2 \times 4$$
  
(C)  $2 \times 2^2$   
(D)  $2 \times 2^3 -$ 

1

$$a^{0} = 1$$
 (Base Case:  $b = 0$ )  
 $a^{b} = a \times a^{b-1}$  (Recursive Case:  $b > 0$ )

Using the recursive definition of exponentiation, what is the value of the following highlighted subexpression?

$$2^3 = 2 \times 2^2 = ?$$

(A)  $2 \times 2$ (B)  $2 \times 2^{1}$ (C) 4(D)  $2 \times 2 \times 2^{1}$ 

$$a^{0} = 1$$
 (Base Case:  $b = 0$ )  
 $a^{b} = a \times a^{b-1}$  (Recursive Case:  $b > 0$ )

Using the recursive definition of exponentiation, what is the value of the following highlighted subexpression?

$$2^3 = 2 \times 2^2 = 2 \times 2 \times 2^1 =?$$

(A)  $2 \times 2^{0}$ (B)  $2 \times 1$ (C) 2(D)  $2 \times 2 \times 2 \times 2^{0}$ 

$$a^{0} = 1$$
 (Base Case:  $b = 0$ )  
 $a^{b} = a \times a^{b-1}$  (Recursive Case:  $b > 0$ )

Using the recursive definition of exponentiation, what is the value of the following highlighted subexpression?

$$2^3 = 2 \times 2^2 = 2 \times 2 \times 2^1 = 2 \times 2 \times 2 \times 2^0 = ?$$

(A)  $2 \times 2^{-1}$ (B)  $2 \times 1$ (C) 2(D) 1

$$a^0 = 1$$
 (Base Case:  $b = 0$ )  
 $a^b = a \times a^{b-1}$  (Recursive Case:  $b > 0$ )

Using the recursive definition of exponentiation, what is the value of the following highlighted subexpression?

$$2^3 = 2 \times 2^2 = 2 \times 2 \times 2^1 = 2 \times 2 \times 2 \times 2^0 = 2 \times 2 \times 2 \times 1$$

(A) 8

(B) 2<sup>3</sup>

(C) Both of the above

(D) None of the above—we don't have a definition for multiplication!

Let's define multiplication recursively. Consider multiplying two natural numbers m and n.

$$m \times ? = ?$$
(Base Case:  $n = ?$ ) $m \times n = ?$ (Recursive Case: ???)

What is the most basic instance of multiplication which requires no recursion?

(A) 
$$n = 1$$
  
(B)  $n = 0$   
(C)  $n = m$   
(D)  $n = m - 1$ 

Let's define multiplication recursively. Consider multiplying two natural numbers m and n.

$$m \times 0 = 0$$
(Base Case:  $n = 0$ ) $m \times n = ???$ (Recursive Case:  $n > 0$ )

How would we move n > 0 closer to the base case?

- (A) Add 1 to n(B) Subtract 1 from n(C) Divide a to 2
- (C) Divide *n* by 2
- (D) Multiply n by m

Let's define multiplication recursively. Consider multiplying two natural numbers m and n.

$$m \times 0 = 0$$
(Base Case:  $n = 0$ ) $m \times n = m \times (n-1) + ???$ (Recursive Case:  $n > 0$ )

So the recursive case must invoke  $m \times (n-1)$ , as that moves us closer to the base case.

What can we add to the value of  $m \times (n-1)$  to get the desired result (i.e., a value equal to the product of m and n)?

(D) 
$$m \times (n-1)$$

$$m \times 0 = 0$$
 (Base Case:  $n = 0$ )  
 $m \times n = m \times (n - 1) + m$  (Recursive Case:  $n > 0$ )

Using the recursive definition of multiplication, what is the value of the following highlighted subexpression?

4 × 3 =?

(A)  $4 \times 3 + 3$ (B)  $4 \times 2 + 4$ (C)  $4 \times 4 + 3$ (D) 12

$$m \times 0 = 0$$
 (Base Case:  $n = 0$ )  
 $m \times n = m \times (n - 1) + m$  (Recursive Case:  $n > 0$ )

Using the recursive definition of multiplication, what is the value of the following highlighted subexpression?

$$4 \times 3 = 4 \times 2 + 4 = ?$$

(A)  $4 \times 1 + 2$ (B)  $4 \times 1 + 4$ (C)  $4 \times 1 + 4 + 4$ (D) 4 + 4

$$m \times 0 = 0$$
 (Base Case:  $n = 0$ )  
 $m \times n = m \times (n - 1) + m$  (Recursive Case:  $n > 0$ )

Using the recursive definition of multiplication, what is the value of the following highlighted subexpression?

$$4 \times 3 = 4 \times 2 + 4 = 4 \times 1 + 4 + 4 =?$$

(A) 4  
(B) 
$$4 \times 0 + 1$$
  
(C)  $4 \times 0 + 4$   
(D)  $4 + 0$ 

$$m \times 0 = 0$$
 (Base Case:  $n = 0$ )  
 $m \times n = m \times (n - 1) + m$  (Recursive Case:  $n > 0$ )

Using the recursive definition of multiplication, what is the value of the following highlighted subexpression?

$$4 \times 3 = 4 \times 2 + 4 = 4 \times 1 + 4 + 4 = 4 \times 0 + 4 + 4 = ?$$

(A) 4  
(B) 
$$4 \times -1 + 4$$
  
(C) 0  
(D)  $4 \times 0$ 

$$m \times 0 = 0$$
 (Base Case:  $n = 0$ )  
 $m \times n = m \times (n - 1) + m$  (Recursive Case:  $n > 0$ )

Using the recursive definition of multiplication, what is the value of the following highlighted subexpression?

 $4 \times 3 = 4 \times 2 + 4 = 4 \times 1 + 4 + 4 = 4 \times 0 + 4 + 4 + 4 = 0 + 4 + 4 + 4 = ?$ 

- (A) 12
- (B) 4 × 3
- (C) Both of the above
- (D) None of the above—we don't have a definition for addition!

#### Digression: Turtles All The Way Down An Apocryphal Tale of Infinite Regress

- Old Lady: The world is really a flate plate supported on the back of a giant turtle.
- Scientist: And what does the turtle stand on?
- Old Lady: Nice try, but it's turtles all the way down!

Do you suppose we could define addition of natural numbers recursively?

- (A) Sure, why not
- (B) No, now you're just pushing it

## Peano Arithmetic

#### Definitions (Axioms)

- $\star 1.$  0 is a natural number
  - 2. For every natural number x, x = x
  - 3. For any two natural numbers x and y, if x = y then y = x
  - 4. For any three natural numbers x, y, and z, if x = y and y = z, then x = z
  - 5. For any objects a and b, if a is a natural number and a = b, then b is a natural number
- $\star 6$ . For every natural number *n*, *n'* is also a natural number
  - 7. For every natural number n, n' = 0 is false
  - 8. For all natural numbers m and n, if m' = n' then m = n
  - 9. Principle of Induction (we'll cover this later!)

Let's define addition recursively.

$$m+?=?$$
(Base Case:  $n=?$ ) $m+?=?$ (Recursive Case:  $n=?$ )

What is the most basic instance of addition which requires no recursion? (A) n = 0'(B) n = 1(C) n = 0(D) n = m

Let's define addition recursively.

$$m + 0 = ?$$
(Base Case:  $n = 0$ ) $m + ? = ?$ (Recursive Case:  $n = ?$ )

What should be the value of m + 0?

(A) 0
(B) m'
(C) m
(D) n

Let's define addition recursively.

$$m + 0 = m$$
(Base Case:  $n = 0$ ) $m + ? = ?$ (Recursive Case:  $n = ?$ )

Intuitively, we want to recurse when n > 0. What does *n* look like when it's not 0?

(A) 
$$n = 0'$$
  
(B)  $n = n'$   
(C)  $n = m'$   
(D)  $n = p'$  for some natural number  $p$ 

Let's define addition recursively.

$$m + 0 = m$$
(Base Case:  $n = 0$ ) $m + p' = ?$ (Recursive Case:  $n = p'$ )

When n = p', which of the following values will be closer to the base case of n = 0?

- (A) n
  (B) n'
  (C) p
- (D) p' 1

Let's define addition recursively.

$$m + 0 = m$$
 (Base Case:  $n = 0$ )  
 $m + p' = \underbrace{m + p}_{?}$  (Recursive Case:  $n = p'$ )

Suppose we recursively invoke addition on m and p by saying m + p. What do we conceptually need to do to the result of m + p in order to get the proper sum, m + p'?

- (A) Add 1 to m + p
- (B) Subtract 1 from m + p
- (C) Multiply m + p by 2
- (D) Divide m + p by 2

Let's define addition recursively.

$$m + 0 = m$$
 (Base Case:  $n = 0$ )  
 $m + p' = \underbrace{m + p}_{?}$  (Recursive Case:  $n = p'$ )

Suppose we recursively invoke addition on m and p by saying m + p. How would we write the value that is 1 more than m + p using the notation of Peano Arithmetic?

(A) 
$$m + p + 0'$$
  
(B)  $(m + p)'$   
(C)  $m + p'$   
(D)  $m' + p$ 

$$m + 0 = m$$
 (Base Case:  $n = 0$ )  
 $m + p' = m' + p$  (Recursive Case:  $n = p'$ )

Using the recursive definition of addition, what is the value of the following highlighted subexpression?

$$2 + 3 = ?$$

(A) 5

(B) 0<sup>'''''</sup>

- (C) 3+1
- (D) None of the above

$$m + 0 = m$$
 (Base Case:  $n = 0$ )  
 $m + p' = m' + p$  (Recursive Case:  $n = p'$ )

Using the recursive definition of addition, what is the value of the following highlighted subexpression?

$$0'' + 0''' = ?$$

(A) 5
(B) 0'''''
(C) 0''' + 0'''
(D) 0' + 0''''

$$m + 0 = m$$
 (Base Case:  $n = 0$ )  
 $m + p' = m' + p$  (Recursive Case:  $n = p'$ )

Using the recursive definition of addition, what is the value of the following highlighted subexpression?

$$0'' + 0''' = 0''' + 0'' = ?$$

(A) 0'''' + 0'(B) 0'''''(C) 0'' + 0''''(D) 0' + 0''''

$$m + 0 = m$$
 (Base Case:  $n = 0$ )  
 $m + p' = m' + p$  (Recursive Case:  $n = p'$ )

Using the recursive definition of addition, what is the value of the following highlighted subexpression?

$$0'' + 0''' = 0''' + 0'' = 0'''' + 0' = ?$$

(A) 0<sup>'''</sup> + 0<sup>''</sup>
(B) 0<sup>'''''</sup>
(C) 0<sup>'''''</sup> + 0
(D) 0<sup>'</sup> + 0<sup>''''</sup>

$$m + 0 = m$$
 (Base Case:  $n = 0$ )  
 $m + p' = m' + p$  (Recursive Case:  $n = p'$ )

Using the recursive definition of addition, what is the value of the following highlighted subexpression?

$$0'' + 0''' = 0''' + 0'' = 0'''' + 0' = 0'''' + 0 =?$$

(A) 0'''''

(B) 5

(C) 
$$0'' + 0'''$$

(D) None of the above

#### Definition (Algorithm)

An algorithm is a step-by-step procedure for accomplishing a task

#### Example (Quicksort)

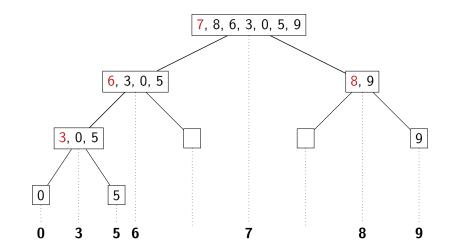
Consider a sequence of elements that can be compared. To sort the sequence recursively, we can do the following:

- If the sequence has < 2 items, it's sorted
- Otherwise,

(Base Case) (Recursive Case)

- Pick a pivot (typically the leftmost element of the sequence)
- Recursively sort the left subsequence of items < the pivot
- Recursively sort the right subsequence of items > the pivot
- Order thus: left subsequence, pivot, right subsequence

#### Quicksort Example



#### **Binary Search**

#### Definition

Suppose you want to search through a sorted sequence,

$$S = [S_1, S_2, S_3, \ldots, S_n]$$

for a particular element, x.

In general, you're always searching between two indices: L and R. Initially, it would be between L = 1 and R = n.

• Let 
$$M = \lfloor (L+R)/2 \rfloor$$

- If L > R, then x is not in S
- If  $S_M = x$ , then x is in S
- If  $S_M > x$ , then search between L and M 1
- If  $S_M < x$ , then search between M + 1 and R

(Base Case)

(Base Case)

(Recursive Case)

(Recursive Case)

Recursion is used to define a lot of things...

- Certain sequences of numbers (e.g., Fibonacci numbers)
- Language (e.g., the definition of a WFF)
- Collections of things (e.g., natural numbers)
- Data types in some programming languages
- The semantics of computable problems

• . . .

So how do we prove facts about recursively-defined things?

#### Definition

Consider a function f upon natural numbers:

$$\sum_{i=s}^{s} f(i) = f(s) \qquad (Base Case: e = s)$$

$$\sum_{i=s}^{e} f(i) = f(e) + \sum_{i=s}^{e-1} f(i) \qquad (Recursive Case: e > s)$$

Let's consider how to prove the following equality is true for any natural number, n.

$$\sum_{i=0}^n i = \frac{n(n+1)}{2}$$

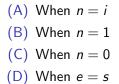
Which proof method do you suppose we need to use here?

- (A) Proof by Example: when *n* matches the base case, we can simply plug it in and see if it works
- (B) Exhaustive Proof: we need to show it holds true for every possible n
- (C) Proof by Cases: depending on whether *n* matches the base case or the recursive case
- (D) Direct Proof: there must be a general way to prove the property

Let's consider how to prove the following equality is true for any natural number, n.

$$\sum_{i=0}^{n} i = \frac{n(n+1)}{2}$$

At what value of *n* would  $\sum_{i=0}^{n} i$  be equal to its base case?



Let's consider how to prove the following equality is true for any natural number, n.

$$\sum_{i=0}^{n} i = \frac{n(n+1)}{2}$$

Consider the case where n = 0.

• 
$$\sum_{i=0}^{n} i =?$$
  
•  $n(n+1)/2 =?$ 

What is the value of the left side of the equality?

(A) 
$$\sum_{i=0}^{0} i$$
  
(B)  $\sum_{i=0}^{n} 0$   
(C)  $\sum_{i=n}^{0} i$   
(D)  $\sum_{i=0}^{n} i$ 

Let's consider how to prove the following equality is true for any natural number, n.

$$\sum_{i=0}^{n} i = \frac{n(n+1)}{2}$$

Consider the case where n = 0.

• 
$$\sum_{i=0}^{n} i = \sum_{i=0}^{0} i = 0$$
  
•  $n(n+1)/2 = ?$ 

What is the value of the right side of the equality?

(A) 
$$0(0+1)/2$$
  
(B) 0

(C) 
$$\sum_{i=0}^{0} i$$

(D) All of the above 
$$(D)$$

Let's consider how to prove the following equality is true for any natural number, n.

$$\sum_{i=0}^{n} i = \frac{n(n+1)}{2}$$

Consider the case where n = 0.

• 
$$\sum_{i=0}^{n} i = \sum_{i=0}^{0} i = 0$$
  
•  $n(n+1)/2 = 0(0+1)/2 = 0$ 

So, the equality holds when the  $\sum$  is at the base case.

What do we know about the value of n when the  $\sum$  uses the recursive case?

(A) 
$$n > e$$
  
(B)  $n > s$ 

(C) 
$$n > 0$$

(D) 
$$n > i$$

Let's consider how to prove the following equality is true for any natural number, n.

$$\sum_{i=0}^{n} i = \frac{n(n+1)}{2}$$

Consider the case where n > 0. What is the value of  $\sum_{i=0}^{n} i$ ?

(A) 
$$\sum_{i=0}^{n-1} i$$
  
(B)  $i + \sum_{i=0}^{n-1} i$   
(C)  $0 + \sum_{i=0}^{n-1} i$   
(D)  $n + \sum_{i=0}^{n-1} i$ 

Let's consider how to prove the following equality is true for any natural number, n.

$$\sum_{i=0}^{n} i = \frac{n(n+1)}{2}$$
$$n + \sum_{i=0}^{n-1} i \stackrel{?}{=} \frac{n(n+1)}{2}$$

Consider the case where n > 0. What is the value of  $\sum_{i=0}^{n-1} i$ ? (A)  $\sum_{i=0}^{n-2} i$ (B)  $(n-1) + \sum_{i=0}^{n-2} i$ (C)  $0 + \sum_{i=0}^{n-2} i$ (D) None of the above